

ELEC 405/511
Error Control Coding

Reed-Solomon Codes

Irving Reed (1923-2012) Gus Solomon (1930-1996)



Polynomial Codes Over Certain Finite Fields, 1960

Reed-Solomon Codes

- Nonbinary BCH codes
- Consider $GF(q)$ ($q=p^r$, p prime)
- To construct a t error correcting nonbinary BCH code with symbols from $GF(q)$, we use the same technique as for binary BCH codes.
- Roots of $g(x)$ are in $GF(q^m)$, $n \mid q^m - 1$
 $n - k \leq 2mt$ product of at most $2t$ minimal polynomials of degree m
 $d \geq 2t + 1$

- Choose $2t$ consecutive powers of α , an element of order n in $\text{GF}(q^m)$, as roots of $g(x)$.
- For RS codes, $m=1$ and α is a primitive element in $\text{GF}(q)$, then

$$n = q-1$$

$$n-k \leq 2t \rightarrow n-k = 2t$$

$$d \geq 2t+1 \rightarrow d \geq n-k+1$$

Singleton Bound

- **Theorem 4-10** Singleton bound

The minimum distance for an (n,k) linear code is bounded by

$$d \leq n - k + 1$$

- For an RS code $d \geq n - k + 1$, so $d = n - k + 1$ and all RS codes meet the Singleton bound
 - they are **optimal** $(n,k,n-k+1)$ codes, $n = q - 1$
- Codes that meet the Singleton bound are called Maximum Distance Separable (MDS)

Reed-Solomon Codes – Minimal Polynomials

- Coefficients of $g(x)$ are in $\text{GF}(q)$, roots of $g(x)$ are also in $\text{GF}(q)$.
- Minimal polynomial of α is $x-\alpha$. There are no conjugates since $\alpha^q = \alpha^{q-1}\alpha = \alpha$.
- BCH: $M_1(x) = (x - \alpha)(x - \alpha^q)(x - \alpha^{q^2}) \dots$
RS: $M_1(x) = (x - \alpha)$
- RS codes are a subclass of BCH codes with $m = 1$.

Example 8-4 $t=2$ GF(8)

- $n = 8-1 = 7$ Form GF(8) from x^3+x+1

$$\alpha^0 \quad 1$$

$$\alpha^1 \quad \alpha$$

$$\alpha^2 \quad \alpha^2$$

$$\alpha^3 \quad \alpha + 1$$

$$\alpha^4 \quad \alpha^2 + \alpha$$

$$\alpha^5 \quad \alpha^2 + \alpha + 1$$

$$\alpha^6 \quad \alpha^2 + 1$$

$$g(x) = (x - \alpha)(x - \alpha^2)(x - \alpha^3)(x - \alpha^4) \\ = x^4 + \alpha^3 x^3 + x^2 + \alpha x + \alpha^3$$

$$\mathbf{H} = \begin{bmatrix} 1 & \alpha & \alpha^2 & \alpha^3 & \alpha^4 & \alpha^5 & \alpha^6 \\ 1 & \alpha^2 & \alpha^4 & \alpha^6 & \alpha & \alpha^3 & \alpha^5 \\ 1 & \alpha^3 & \alpha^6 & \alpha^2 & \alpha^5 & \alpha & \alpha^4 \\ 1 & \alpha^4 & \alpha^8 & \alpha^5 & \alpha^2 & \alpha^6 & \alpha^3 \end{bmatrix}$$

- (7,3,5) RS code

Comparison: RS vs Binary BCH

- RS: $n = q^m - 1$ $q = 8, m = 1$ (7,3,5)

$$g(x) = (x - \alpha)(x - \alpha^2)(x - \alpha^3)(x - \alpha^4)$$

- Binary BCH: $n = q^m - 1$ $q = 2, m = 3$ (7,1,7)

$$g(x) = (x - \alpha)(x - \alpha^2)(x - \alpha^3)(x - \alpha^4)(x - \alpha^6)(x - \alpha^5)$$

- RS code: $q^k = 8^3 = 512$ codewords
- Binary BCH code: $q^k = 2^1 = 2$ codewords

Comparison: RS vs Binary BCH

- (7,3,5) RS code: $8^3 = 512$ codewords = 2^9
- Each symbol can be represented as 3 bits, a codeword has $n = 7$ symbols = 21 bits and $k = 3$ data symbols = 9 bits.
- The (7,3,5) RS code can be considered as a (21,9) binary code.
- $t = 2$ symbols - since 5 bit errors may cover 3 symbols, corrects any burst error of 4 bits or less.

Example 8-5 $t=3$ GF(64)

- $n = 64-1 = 63$
- α a root of the primitive polynomial x^6+x+1
$$g(x) = (x - \alpha)(x - \alpha^2)(x - \alpha^3)(x - \alpha^4)(x - \alpha^5)(x - \alpha^6)$$
$$= x^6 + \alpha^{59}x^5 + \alpha^{48}x^4 + \alpha^{43}x^3 + \alpha^{55}x^2 + \alpha^{10}x + \alpha^{21}$$
- $(63,57,7)$ RS code
- $64^{57} = 8.96 \times 10^{102}$ codewords
- $64^{63} = 6.16 \times 10^{113}$ vectors
- sphere volume is 9.94×10^9 so the spheres fill about 14.5% of the vector space

GF(7) Example

- RS codes can be constructed over any finite field
- Consider $q = 7$ so that $n = q-1 = 6$, and $t = 2$
- First find a primitive element in GF(7)

$\phi(6) = 2$ so two primitive elements

$3^1=3 \quad 3^2=2 \quad 3^3=6 \quad 3^4=4 \quad 3^5=5 \quad 3^6=1 \rightarrow 3$ is primitive

$$\begin{aligned} b=1 \quad g(x) &= (x-3^1)(x-3^2)(x-3^3)(x-3^4) \\ &= (x-3)(x-2)(x-6)(x-4) \quad (6,2,5) \text{ RS code} \end{aligned}$$

$$\begin{aligned} b=2 \quad g(x) &= (x-3^2)(x-3^3)(x-3^4)(x-3^5) \\ &= (x-2)(x-6)(x-4)(x-5) \quad (6,2,5) \text{ RS code} \end{aligned}$$

- One can pick any group of consecutive roots

$$g(x) = (x-3^1)(x-3^2)(x-3^3)$$

$$= (x-3)(x-2)(x-6) \quad (6,3,4) \text{ RS code}$$

$$= x^3+3x^2+x+6$$

$$g(x) = (x-3^2)(x-3^3)(x-3^4)$$

$$= (x-2)(x-6)(x-4) \quad (6,3,4) \text{ RS code}$$

$$= x^3+2x^2+2x+1 = g^*(x) \quad \text{self reciprocal}$$

$$g(x) = (x-3^1)(x-3^2)(x-3^3)(x-3^4)(x-3^5)$$

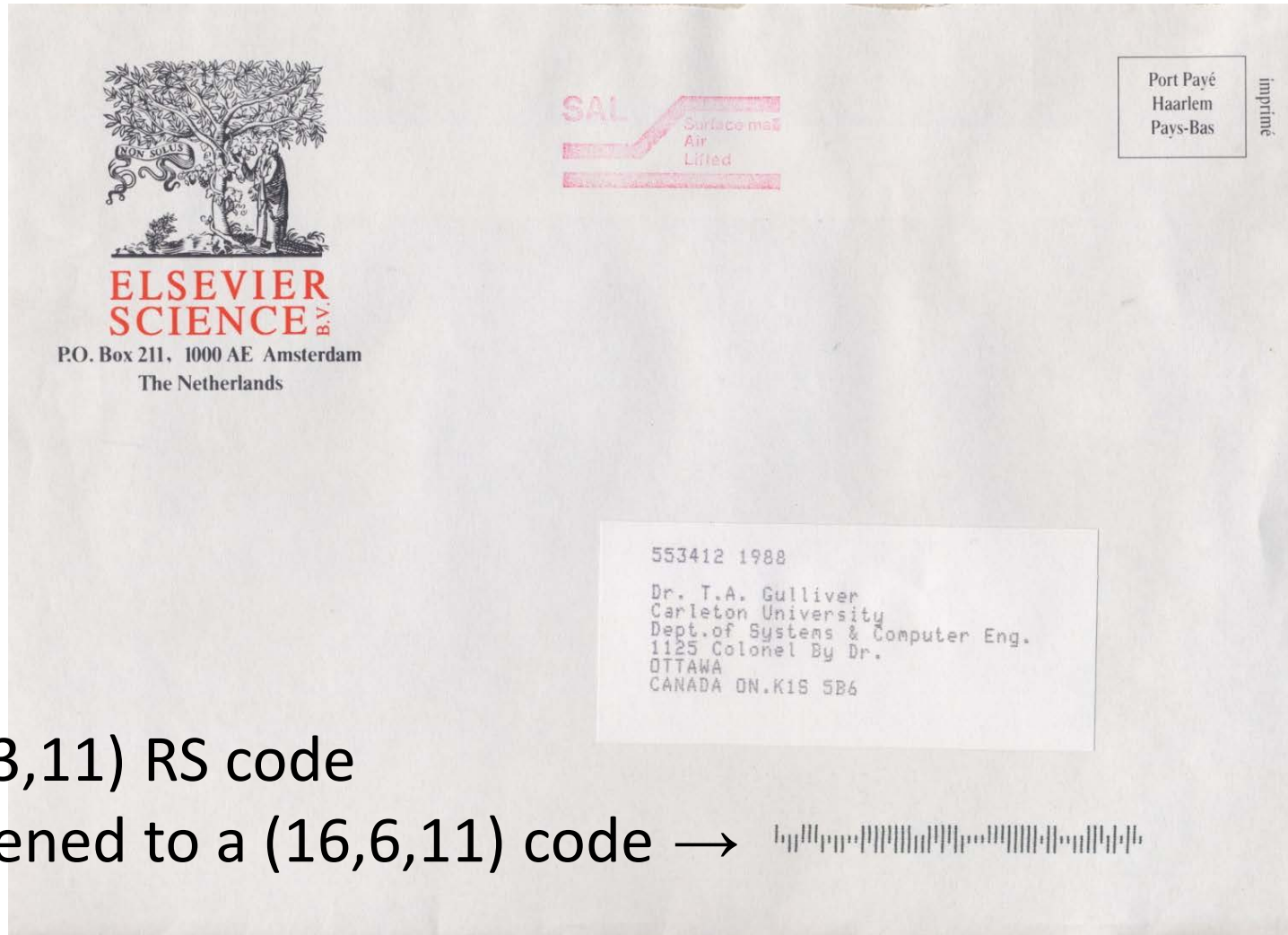
$$= (x-3)(x-2)(x-6)(x-4)(x-5) \quad (6,1,6) \text{ RS code}$$

$$= x^5+x^4+x^3+x^2+x+1 = g^*(x) \quad \text{self reciprocal}$$

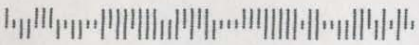
Properties of RS Codes

- The dual code of an RS code is also MDS
 - C (6,2,5) code over GF(7)
 - C^\perp (6,4,3) code over GF(7)
- Since RS codes are cyclic codes, they can always be put in systematic form $x^{n-k}m(x)+d(x)$
- A shortened RS codes is MDS
$$(n,k,n-k+1) \rightarrow (n-u,k-u,n-k+1) \quad (6,4,3) \rightarrow (5,3,3)$$
- A punctured RS code is MDS
$$(n,k,n-k+1) \rightarrow (n-u,k,n-k-u+1) \quad (6,4,3) \rightarrow (5,4,2)$$

Example: Bar Codes over GF(64)



(63,53,11) RS code

shortened to a (16,6,11) code → 

Extended RS Codes

- An (n,k) RS code over $GF(q)$ with $n = q-1$ can be extended twice to a $(q+1,k)$ MDS code
- There is a technique for constructing such codes which are cyclic
- A very few RS codes can be triply extended to obtain an MDS code
 - $k = 3$ or $n-k = 3$ and $q = 2^m$
 - In this case $n = q+2$

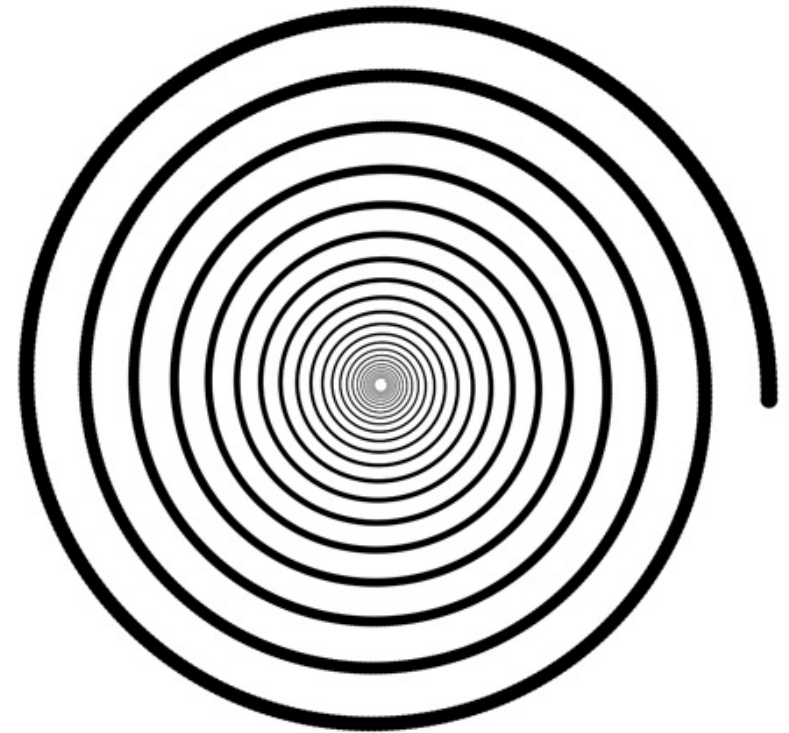
Example: NASA/JPL Code

- $q = 256, n = q-1 = 255$
- $(255,223,33)$ RS code over $GF(2^8)$

$$\frac{\# \text{ of codewords} \times \text{volume}}{\text{size of vector space}} = 2.78 \times 10^{-14}$$

Example: Compact Discs

- 44.1 kHz sample rate
- 16 bit stereo samples
- $2 \times 16 \times 44100 = 1.41$ Mbps
- Original CD capacity: 74 minutes of audio or 650 MB of data
- Data stored on a spiral, not concentric circles
 - length 5.38 km
 - velocity 1.2 m/s



Kees Schouhamer Immink (1946 -)



Sources of Error

- 1) Defects caused during disc production
 - inferior disc pits and bubbles during disc formation
 - defects in the aluminum film and a poor reflective index
- 2) Defects caused in handling
 - fingerprints and scratches
 - dust
- 3) Variations and disturbances during playback
 - disturbance of the servo mechanism
- 4) Jitter - time variation of the signal
- 5) Interference

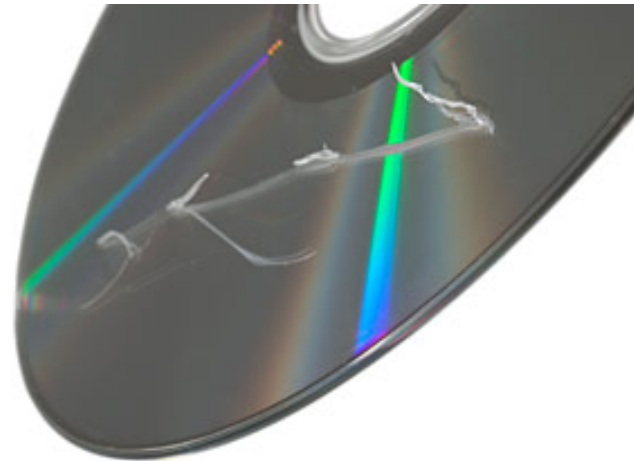
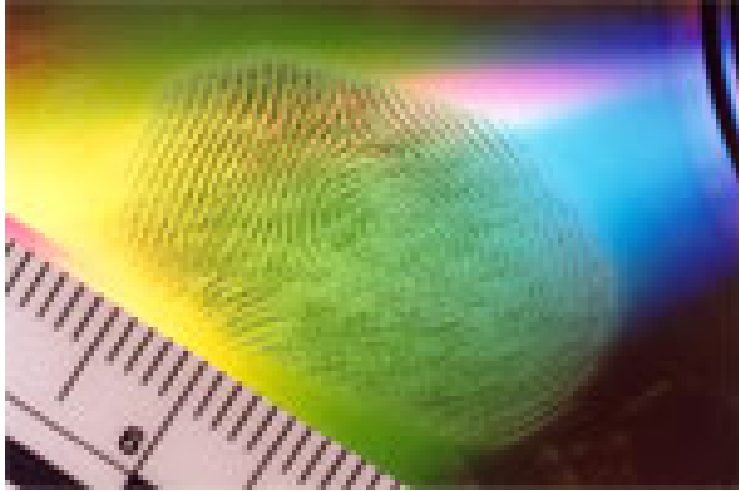
(1)-(3) cause burst errors

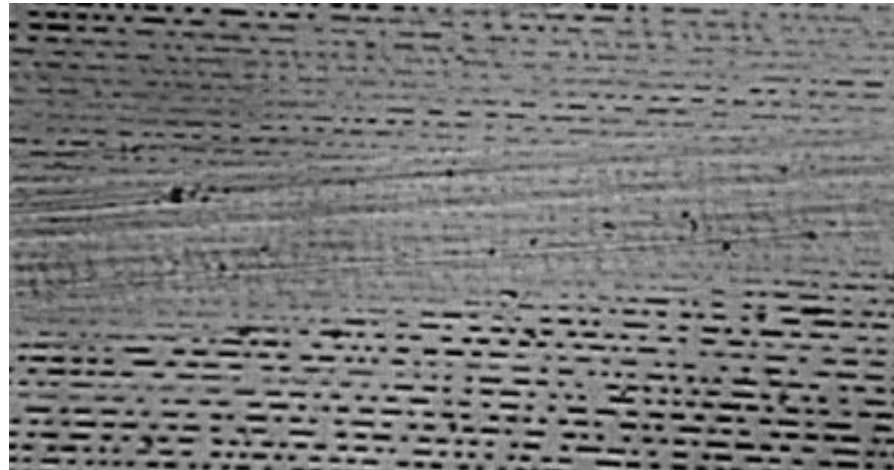
(4) and (5) cause random errors

Causes of Disc Errors

- Fingerprints cause 43% of errors
- General wear and tear causes 25% of errors
- Player-related issues cause 15% of errors
- User-related issues cause 12% of errors
- Manufacturer errors cause 2% of errors

Causes of Disc Errors





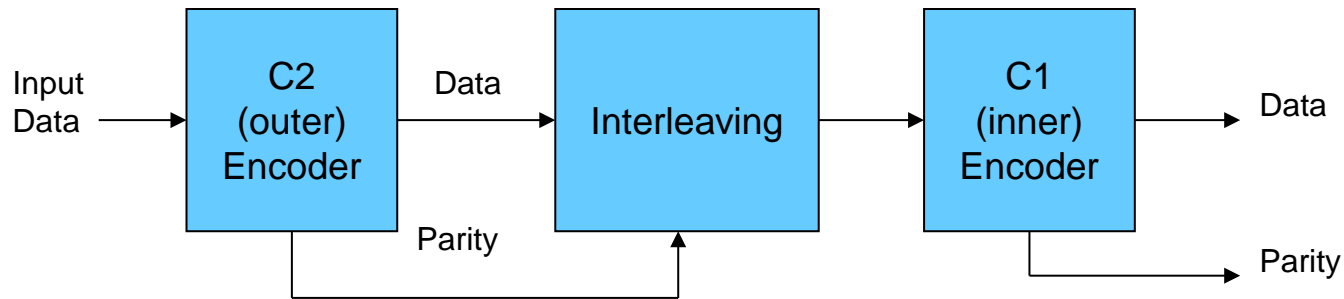
Error Correction

- Reed-Solomon code
 - (255,251,5) code over $GF(2^8)$
- Shortened to a (28,24,5) outer code
- These codewords are interleaved to reduce the effects of burst errors
- (32,28,5) inner code
- Overall code rate is

$$\frac{24}{28} \times \frac{28}{32} = 0.75$$

CIRC Encoder

- CIRC – Cross Interleaved Reed-Solomon Code



- Interleaving disperses the codewords so they are not contiguous on the disc
- mitigates long burst errors associated with scratches and fingerprints
 - Maximum correctable burst error length
 - 4000 bits = 2.5 mm

Encoding Algorithm

- Samples are split into two 8 bit symbols
- Six samples from each channel are grouped to obtain 24 symbols
- Four outer RS code parity symbols are generated to give a frame of 28 symbols
- Symbols are interleaved over 109 frames
- Four inner RS parity symbols are generated to give 32 symbols
- These frames are also interleaved

Control and Error Correction

- Skips are caused by physical disturbances
 - Wait for disturbance to subside
 - Retry
- Read errors caused by disc/servo problems
 - Detect error
 - Choose location for retry
 - Retry, if it fails interpolate if applicable

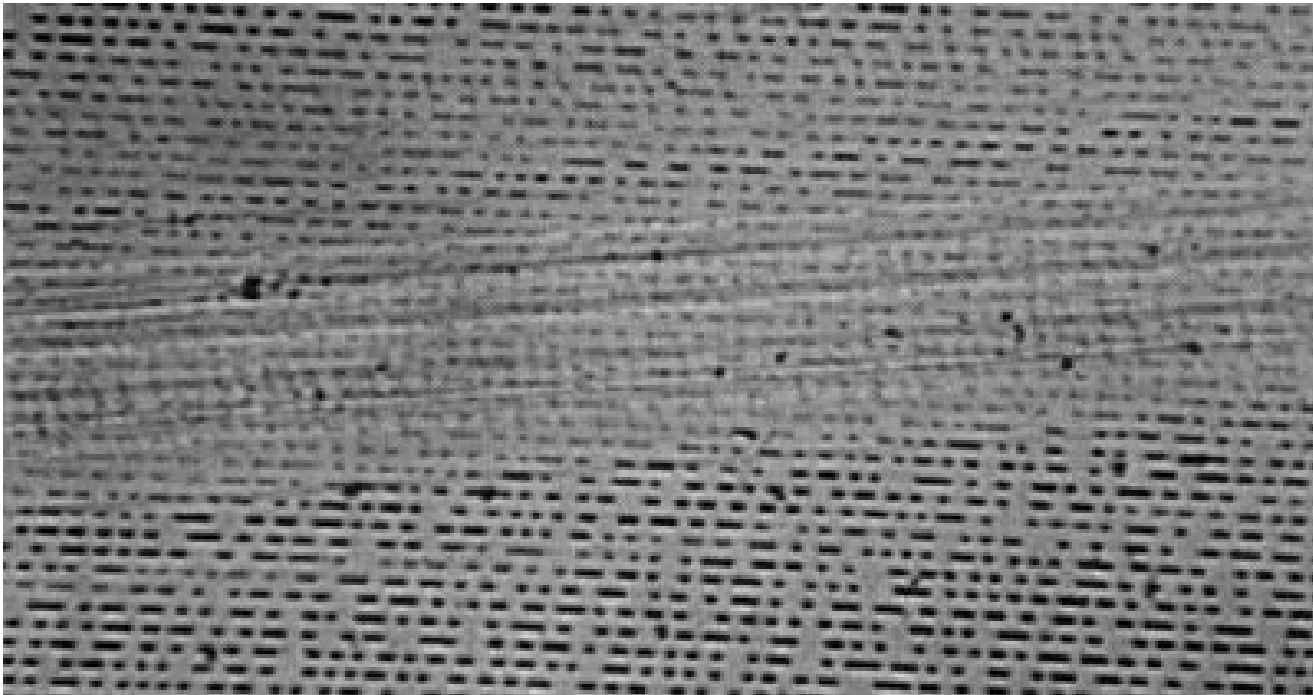
Interpolation

- Used when decoding fails
- Fill missing audio data using adjacent data
 - time or channel
- Only valid for audio CDs

Decoding RS Codes

1. Compute the syndromes
2. Determine the error locator polynomial $\Lambda(x)$
3. Determine the error magnitudes from $\Lambda'(x)$ and $\Omega(x)$
$$\Omega(x) = [1 + S(x)]\Lambda(x)$$
4. Evaluate the error locations and the error values at those locations.

CD Errors due to a Ball Point Pen



A Highly Corroded Disc

- Two minutes can still be played.



Audio Data Format

